

Colorization

CVFX @ NTHU

3 March 2015

Outline

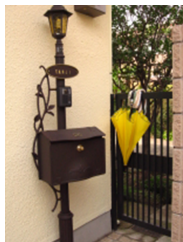
Colorization

The paper

Colorization using optimization

- ▶ Levin, Lischinski, and Weiss
- ▶ SIGGRAPH 2004
- ▶ Useful skills of linear algebra

Example



Advantages

Solving a sparse linear system

- ▶ Many numerical methods can be used

No segmentation required

- ▶ Users only need to draw some color scribbles

Intensity contrast is maintained

BlackMagic/TimeBrush Real-Life-Color

<http://www.blackmagic-color.com>

<http://www.timebrush.com/>

<http://www.blackmagic-color.com/>

<http://www.timebrush.com/>

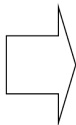
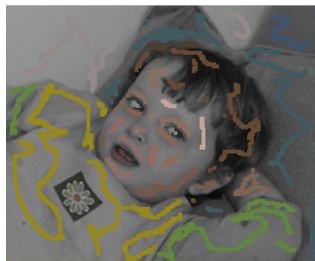
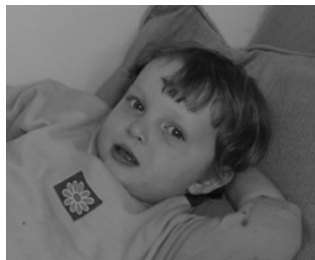


http://www.blackmagic-color.com/alt/bm_samples2.html



http://www.blackmagic-color.com/alt/bm_samples3.html

Colorization



Recoloring



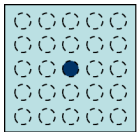
Modification for recoloring

```
sgI = rgb2ntsc(gI);
scI = rgb2ntsc(cI);
ntscIm(:,:,1) = sgI(:,:,1);
ntscIm(:,:,2) = scI(:,:,2);
ntscIm(:,:,3) = scI(:,:,3);

idx = find(sum(cI,3)>=3); % For white scribbles,
tmp1 = ntscIm(:,:,2); % we have to copy from
tmp2 = sgI(:,:,2); % the original image
tmp1(idx) = tmp2(idx); % those colors of the
ntscIm(:,:,2) = tmp1; % corresponding pixels.
tmp1 = ntscIm(:,:,3);
tmp2 = sgI(:,:,3);
tmp1(idx) = tmp2(idx);
ntscIm(:,:,3) = tmp1;
```

Local linearity

Assume that the color (U or V) at a pixel is a linear function of the intensity, and the linear coefficients are the same for all pixels in a small neighborhood


$$U_i \simeq a_k Y_i + b_k$$

Ω_k

(a_k, b_k)

Cost function

The squared error of linear approximation for the whole image

$$J(U) = \sum_k \left(\min_{a_k, b_k} \sum_{i \in \Omega_k} (U_i - a_k Y_i - b_k)^2 \right)$$

Scirbbles?

Constraints

$$U_i = u_i$$

Optimization

Find U that minimizes

$$J(U) = \sum_k \left(\min_{a_k, b_k} \sum_{i \in \Omega_k} (U_i - a_k Y_i - b_k)^2 \right)$$

subject to the constraints

$$U_i = u_i$$

Rewrite the cost function

$$J(U) = \sum_k \left(\min_{a_k, b_k} \sum_{i \in \Omega_k} (U_i - a_k Y_i - b_k)^2 \right)$$

$$J(U) = \sum_k \min_{a_k, b_k} \left\| \begin{pmatrix} Y_1 & 1 \\ Y_2 & 1 \\ \vdots & \vdots \\ Y_{|\Omega_k|} & 1 \end{pmatrix} \begin{pmatrix} a_k \\ b_k \end{pmatrix} - \begin{pmatrix} U_1 \\ U_2 \\ \vdots \\ U_{|\Omega_k|} \end{pmatrix} \right\|^2$$

Let $P_k = \begin{pmatrix} Y_1 & 1 \\ Y_2 & 1 \\ \vdots & \vdots \\ Y_{|\Omega_k|} & 1 \end{pmatrix}$ and $\tilde{U}_k = \begin{pmatrix} U_1 \\ U_2 \\ \vdots \\ U_{|\Omega_k|} \end{pmatrix}$

Optimization (rewritten)

Find U that minimizes

$$J(U) = \sum_k \left(\min_{a_k, b_k} \left\| P_k \begin{pmatrix} a_k \\ b_k \end{pmatrix} - \widetilde{U}_k \right\|^2 \right)$$

subject to the constraints

$$U_i = u_i$$

Subproblem

to minimize

$$Q_k(a_k, b_k) = \left\| P_k \begin{pmatrix} a_k \\ b_k \end{pmatrix} - \widetilde{U}_k \right\|^2$$

is equivalent to find the solution of

$$P_k^T P_k \begin{pmatrix} a_k \\ b_k \end{pmatrix} = P_k^T \widetilde{U}_k$$

the least squares solution

$$\begin{pmatrix} a_k^* \\ b_k^* \end{pmatrix} = (P_k^T P_k)^{-1} P_k^T \widetilde{U}_k$$

Substitute the solution into the sub-cost

$$\begin{pmatrix} a_k^* \\ b_k^* \end{pmatrix} = (P_k^T P_k)^{-1} P_k^T \tilde{U}_k$$

$$(AB)^T = B^T A^T$$

$$Q_k^*(a_k^*, b_k^*) = \left\| P_k (P_k^T P_k)^{-1} P_k^T \tilde{U}_k - \tilde{U}_k \right\|^2$$

$$\|g\|^2 = g^T g$$

$$Q_k^*(a_k^*, b_k^*) = \left\| \underbrace{\left(P_k (P_k^T P_k)^{-1} P_k^T - I \right)}_g \tilde{U}_k \right\|^2$$

$$\|g\|^2 = g^T g$$

$$Q_k^*(a_k^*, b_k^*) = \tilde{U}_k^T \underbrace{\left(P_k (P_k^T P_k)^{-1} P_k^T - I \right)}_g^T \underbrace{\left(P_k (P_k^T P_k)^{-1} P_k^T - I \right)}_g \tilde{U}_k$$

$$\begin{aligned} & P_k (P_k^T P_k)^{-1} P_k^T P_k (P_k^T P_k)^{-1} P_k^T - 2P_k (P_k^T P_k)^{-1} P_k^T + I \\ &= \underbrace{I - P_k (P_k^T P_k)^{-1} P_k^T} \end{aligned}$$

Rewrite the sub-cost

$$\begin{aligned} Q_k^*(a_k^*, b_k^*) &= \tilde{U}_k^T \left(P_k (P_k^T P_k)^{-1} P_k^T - I \right)^T \left(P_k (P_k^T P_k)^{-1} P_k^T - I \right) \tilde{U}_k \\ &= \tilde{U}_k^T \left(I - P_k (P_k^T P_k)^{-1} P_k^T \right) \tilde{U}_k \\ &= \underbrace{\tilde{U}_k^T L_k \tilde{U}_k}_{\left(\quad \right) \square \left(\quad \right)} \end{aligned}$$

$L_k = \left(I - P_k (P_k^T P_k)^{-1} P_k^T \right)$ is a matrix of size $|\Omega|_k$ by $|\Omega|_k$

Details

The (i, j) element of L_k is

$$\begin{aligned} & \left(I - P_k (P_k^T P_k)^{-1} P_k^T \right)_{ij} \\ &= \delta_{ij} - (Y_i \ 1) \begin{pmatrix} \sum_l^{|\Omega_k|} Y_l^2 & \sum_l^{|\Omega_k|} Y_l \\ \sum_l^{|\Omega_k|} Y_l & |\Omega_k| \end{pmatrix}^{-1} \begin{pmatrix} Y_j \\ 1 \end{pmatrix} \end{aligned}$$



$$\delta_{ij} = \begin{cases} 1, & i=j \\ 0, & i \neq j \end{cases}$$

$$I = \begin{pmatrix} 1 & & & \\ & 1 & & \\ & & \dots & \\ & & & 1 \end{pmatrix}$$

$$\begin{pmatrix} Y_1 & 1 \\ Y_2 & 1 \\ \vdots & \\ Y_{|\Omega_k|} & 1 \end{pmatrix} \left\{ \begin{pmatrix} Y_1 & Y_2 & \dots & Y_{|\Omega_k|} \\ 1 & 1 & \dots & 1 \end{pmatrix} \begin{pmatrix} Y_1 & 1 \\ Y_2 & 1 \\ \vdots & \\ Y_{|\Omega_k|} & 1 \end{pmatrix} \right\}^{-1} \begin{pmatrix} Y_1 & Y_2 & \dots & Y_{|\Omega_k|} \\ 1 & 1 & \dots & 1 \end{pmatrix}$$



Compute the inverse

$$\begin{aligned} \begin{pmatrix} \sum_l^{|\Omega_k|} Y_l^2 & \sum_l^{|\Omega_k|} Y_l \\ \sum_l^{|\Omega_k|} Y_l & |\Omega_k| \end{pmatrix}^{-1} &= \frac{\begin{pmatrix} |\Omega_k| & -\sum_l^{|\Omega_k|} Y_l \\ -\sum_l^{|\Omega_k|} Y_l & \sum_l^{|\Omega_k|} Y_l^2 \end{pmatrix}}{|\Omega_k| \sum_l^{|\Omega_k|} Y_l^2 - (\sum_l^{|\Omega_k|} Y_l)^2} \\ &= \frac{|\Omega_k| \begin{pmatrix} 1 & -\mu_k \\ -\mu_k & \sum_l^{|\Omega_k|} Y_l^2 / |\Omega_k| \end{pmatrix}}{|\Omega_k|^2 \sigma_k^2} \\ &= \frac{1}{|\Omega_k| \sigma_k^2} \begin{pmatrix} 1 & -\mu_k \\ -\mu_k & \sum_l^{|\Omega_k|} Y_l^2 / |\Omega_k| \end{pmatrix} \end{aligned}$$

Some efforts

The (i, j) element of L_k

$$\begin{aligned}(L_k)_{ij} &= \left(I - P_k (P_k^T P_k)^{-1} P_k^T \right)_{ij} \\ &= \delta_{ij} - (Y_i \ 1) \begin{pmatrix} \sum_l^{|\Omega_k|} Y_l^2 & \sum_l^{|\Omega_k|} Y_l \\ \sum_l^{|\Omega_k|} Y_l & |\Omega_k| \end{pmatrix}^{-1} \begin{pmatrix} Y_j \\ 1 \end{pmatrix} \\ &= \delta_{ij} - (Y_i \ 1) \frac{1}{|\Omega_k| \sigma_k^2} \begin{pmatrix} 1 & -\mu_k \\ -\mu_k & \sum_l^{|\Omega_k|} Y_l^2 / |\Omega_k| \end{pmatrix} \begin{pmatrix} Y_j \\ 1 \end{pmatrix} \\ &= \delta_{ij} - \frac{1}{|\Omega_k| \sigma_k^2} \left(Y_i Y_j - Y_i \mu_k - Y_j \mu_k + \frac{\sum_l^{|\Omega_k|} Y_l^2}{|\Omega_k|} \right) \\ &= \delta_{ij} - \frac{1}{|\Omega_k| \sigma_k^2} \left(Y_i Y_j - Y_i \mu_k - Y_j \mu_k + \mu_k^2 + \frac{\sum_l^{|\Omega_k|} Y_l^2}{|\Omega_k|} - \mu_k^2 \right) \\ &= \delta_{ij} - \frac{1}{|\Omega_k| \sigma_k^2} \left((Y_i - \mu_k)(Y_j - \mu_k) + \sigma_k^2 \right) \\ &= \delta_{ij} - \frac{1}{|\Omega_k|} \left(1 + \frac{1}{\sigma_k^2} (Y_i - \mu_k)(Y_j - \mu_k) \right)\end{aligned}$$

$$\tilde{U}_k^T L_k \tilde{U}_k$$

Optimization

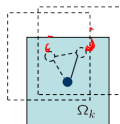
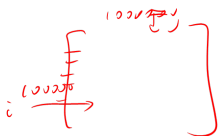
Find U that minimizes

$$J(U) = \sum_k \tilde{U}_k^T L_k \tilde{U}_k = U^T L U$$

subject to the constraints

Solve a sparse linear system

$$\tilde{U}_k = \begin{pmatrix} U_1 \\ U_2 \\ \vdots \\ U_{|\Omega_k|} \end{pmatrix}$$



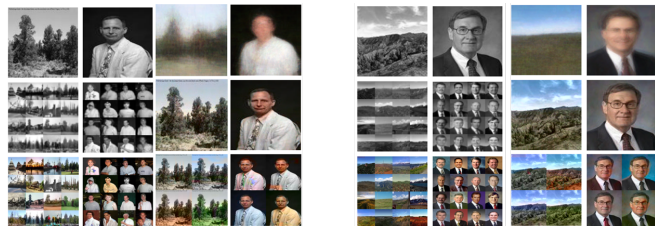
L is a large sparse N -by- N matrix whose (i, j) element is

$$\sum_{k|(i,j) \in \Omega_k} \left(\delta_{ij} - \frac{1}{|\Omega_k|} \left(1 + \frac{1}{\sigma_k^2} (Y_i - \mu_k)(Y_j - \mu_k) \right) \right)$$

N is the number of pixels in the image

Learning from examples

- ▶ 80 million tiny images (Torralba et al.)



Conclusion

Scribbles

- ▶ Cannot fully rely on computers
- ▶ To get good results on hard problems, we must make some efforts too
- ▶ Human-computer interaction, distributed human computing, Amazon MT

Optimization

- ▶ Formulate your problem based on reasonable assumptions
- ▶ You are lucky if the cost function is quadratic and the constraints are linear

Another approach to scribble-based colorization

Fast Image and Video Colorization Using Chrominance Blending

- ▶ Yatziv and Sapiro
- ▶ IEEE Transactions on Image Processing, vol. 15, no. 5, 2006

YCbCr color spaces

- ▶ Y channel: luminance
- ▶ Cb and Cr channels: chrominance

JPEG-YCbCr from 8-bit RGB

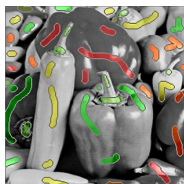
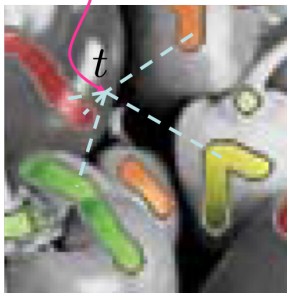
$$\begin{aligned} & \text{=====} \\ Y &= \quad \quad + 0.299 * R \quad + \quad 0.587 * G + \quad 0.114 * B \\ Cb &= 128 - 0.168736 * R - 0.331264 * G + \quad 0.5 * B \\ Cr &= 128 + \quad 0.5 * R - 0.418688 * G - 0.081312 * B \end{aligned}$$

R, G, B in {0, 1, 2, ..., 255}

Idea

chrominance unknown

?



Ω : all pixels

Ω_c : scribbled pixels

Problem

Given

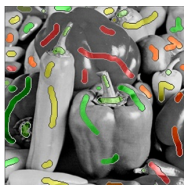
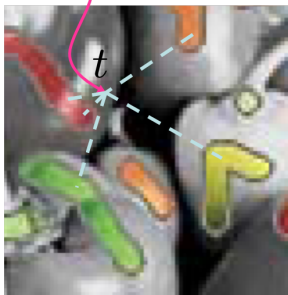
- ▶ the chrominance channels in Ω_c
- ▶ the luminance channel in Ω

To estimate

- ▶ the chrominance channels in $\Omega \setminus \Omega_c$

Chrominance blending

blending



Ω : all pixels

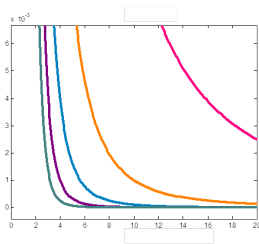
Ω_c : scribbled pixels

How to blend

$$\text{chrominance}(t) \leftarrow \frac{\sum_{\forall c \in \text{chrominances}(\Omega_c)} W(d_c(t))c}{\sum_{\forall c \in \text{chrominances}(\Omega_c)} W(d_c(t))}$$

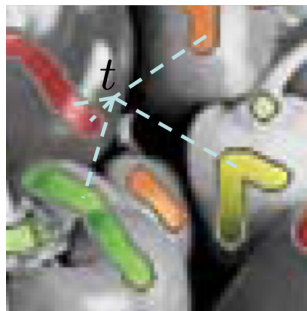
$d_c(t)$: the distance from pixel t to the 'nearest' pixel with chrominance c

$$W(d) = d^{-b}, 1 \leq b \leq 6$$

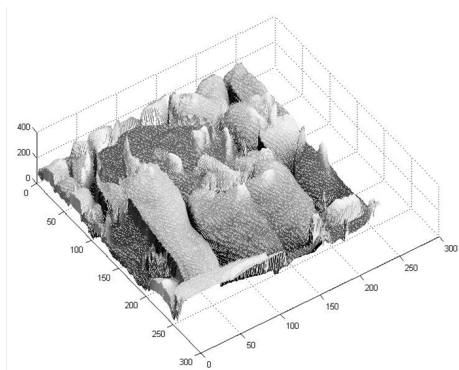
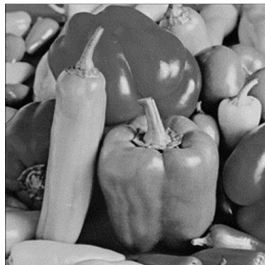


How to define nearness?

$d_c(t)$: the distance from pixel t to the 'nearest' pixel with chrominance c

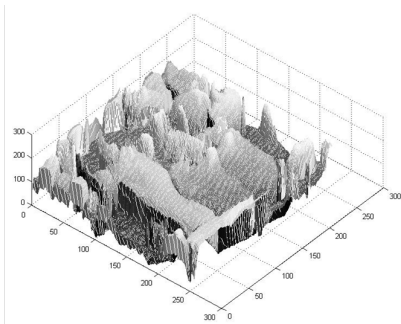
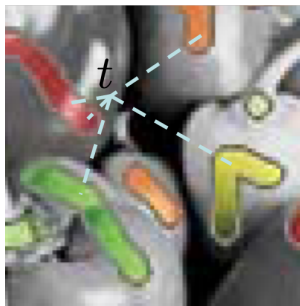


Incorporate luminance information



Incorporate luminance information

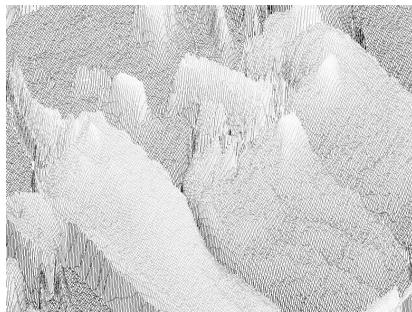
How far do we need to travel (uphill and downhill) from t to a scribbled pixel?



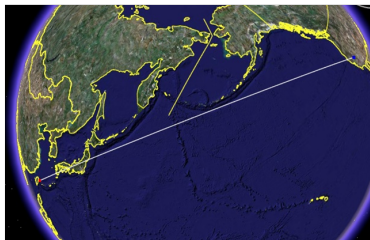
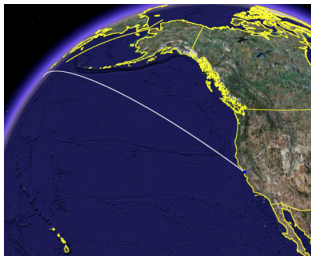
Reasonable?

Consider

- ▶ smooth areas
- ▶ rugged areas
- ▶ edges

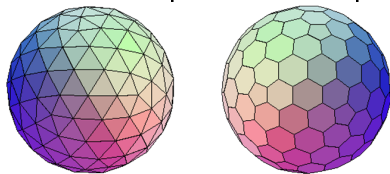


Geodesic distance



Geodesic distance

the shortest path between points on the space



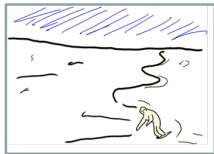
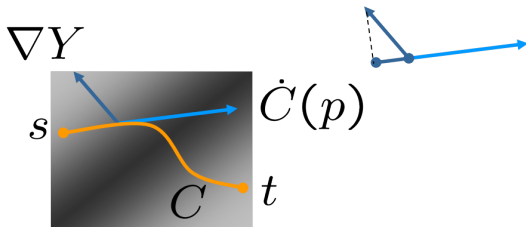
http://en.wikipedia.org/wiki/Geodesic_dome



Definition

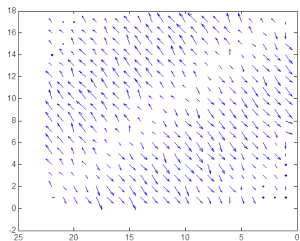
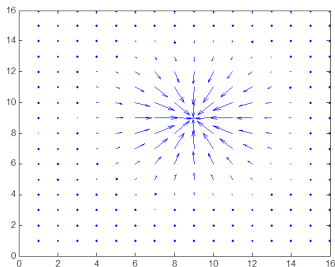
Let s and t be two points in Ω and let $C(p) : [0, 1] \rightarrow \Omega$ be a curve in Ω . Also, let $C_{s,t}$ represent a curve connecting s and t such that $C(0) = s$ and $C(1) = t$. We define the geodesic distance between s and t by

$$d(s, t) \equiv \min_{C_{s,t}} \int_0^1 |\nabla Y \cdot \dot{C}(p)| dp$$



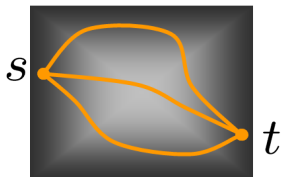
Visualization

$$d(s, t) = \min_{C_{s,t}} \int_0^1 |\nabla Y \cdot \dot{C}(p)| dp$$



Find the shortest path

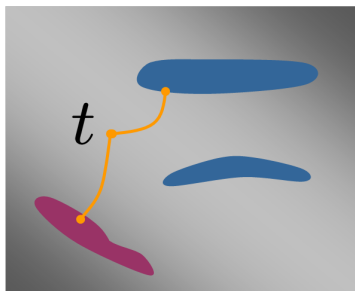
$$d(s, t) = \min_{C_{s,t}} \int_0^1 |\nabla Y \cdot \dot{C}(p)| dp$$



find the 'flatest' path connecting the two points

Distance to a certain chrominance

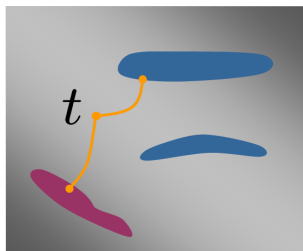
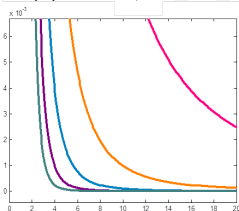
$$d_c(t) := \min_{\forall s \in \Omega_c: \text{chrominance}(s) = c} d(s, t)$$



Chrominance blending

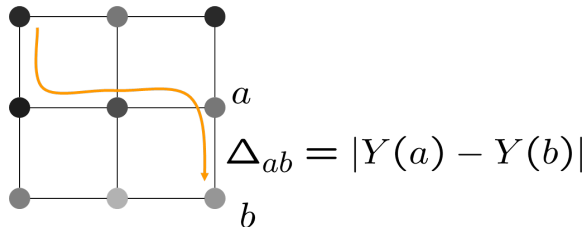
$$\text{chrominance}(t) \leftarrow \frac{\sum_{\forall c \in \text{chrominances}(\Omega_c)} W(d_c(t))c}{\sum_{\forall c \in \text{chrominances}(\Omega_c)} W(d_c(t))}$$

$$W(d) = d^{-b}, 1 \leq b \leq 6$$

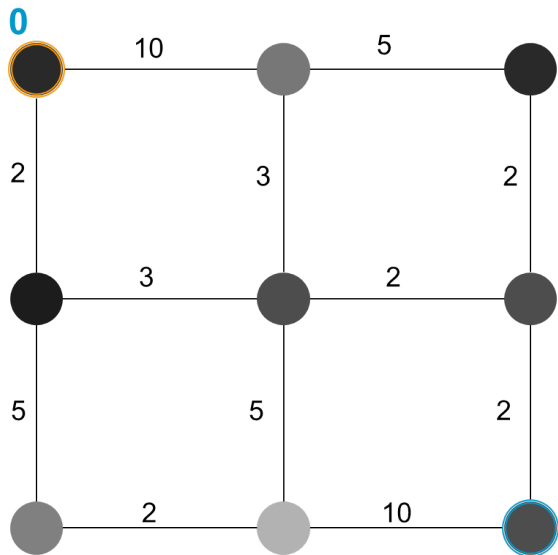


How to compute geodesic distance on an image?

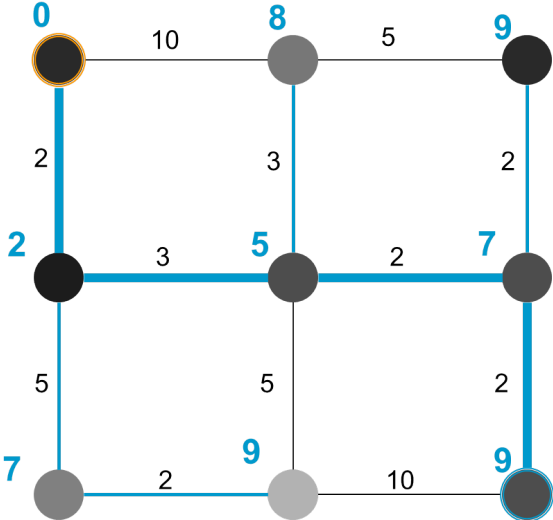
- ▶ Create a graph, and find the shortest path
- ▶ Dijkstra's algorithm
- ▶ Using a priority queue
- ▶ $O((E + V) \log V)$



Dijkstra example



Result



Comparison 1

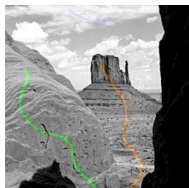


[Yatziv & Sapiro]



[Levin *et al.*]

Comparison 2

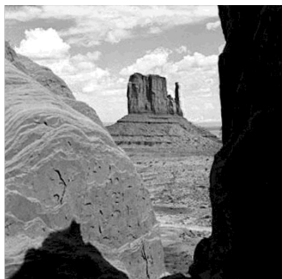
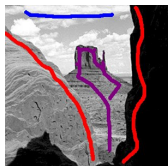


[Yatviz & Sapiro]



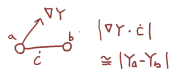
[Levin *et al.*]

Another application: segmentation



Numerical issues

- ▶ Dijkstra's algorithm has a bias when measuring distances on a grid



- ▶ How to compute the gradients more accurately?
- ▶ How to achieve sub-grid accuracy?
- ▶ Level-set methods

Conclusion

Scribbles

- ▶ Interactive colorization tool

Fast blending

- ▶ Geodesic distance
- ▶ Dijkstra's algorithm